



Lecture 8

Counting sort, Bucket sort, Radix sort

CS 161 Design and Analysis of Algorithms

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Counting sort

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 - ▶ If x only appears once in A , then x should go in in $B[j]$.
 - ▶ If x appears more than once in A and we want a stable sort:
 - ▶ Last occurrence of x in A should go in $B[j]$
 - ▶ Next-to-last occurrence of x should go in $B[j - 1]$
 - ▶ etc.

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- ▶ Use an auxiliary array $locator[1..k]$

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- ▶ $locator[x]$ contains the index of the position in the output array B where a key of x should be stored.

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 - ▶ We make several passes over the data to set the values in the locator array before we do the actual sort.

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 - ▶ At the start of the final (sorting) pass, $locator[x]$ contains the number of elements $\leq x$

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- ▶ On the final pass:

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 - ▶ We make several passes over the data to set the values in the locator array before we do the actual sort.
 - ▶ At the start of the final (sorting) pass, $locator[x]$ contains the number of elements $\leq x$
- ▶ On the final pass:
 - ▶ Process the input array A from right to left (!). This makes the counting sort a stable sorting algorithm.

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 - ▶ Output array is $B[1..n]$
- ▶ Use an auxiliary array $locator[1..k]$
- ▶ $locator[x]$ contains the index of the position in the output array B where a key of x should be stored.
 - ▶ We make several passes over the data to set the values in the locator array before we do the actual sort.
 - ▶ At the start of the final (sorting) pass, $locator[x]$ contains the number of elements $\leq x$
- ▶ On the final pass:
 - ▶ Process the input array A from right to left (!). This makes the counting sort a stable sorting algorithm.
 - ▶ When a value of x is encountered in the input array A :
 - ▶ Copy the value into location $locator[x]$ in the output array. That is, store it in location $B[locator[x]]$

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 - ▶ Each integer is in the range $1..k$
 - ▶ Output array is $B[1..n]$
- ▶ Use an auxiliary array $locator[1..k]$
- ▶ $locator[x]$ contains the index of the position in the output array B where a key of x should be stored.
 - ▶ We make several passes over the data to set the values in the locator array before we do the actual sort.
 - ▶ At the start of the final (sorting) pass, $locator[x]$ contains the number of elements $\leq x$
- ▶ On the final pass:
 - ▶ Process the input array A from right to left (!). This makes the counting sort a stable sorting algorithm.
 - ▶ When a value of x is encountered in the input array A :
 - ▶ Copy the value into location $locator[x]$ in the output array. That is, store it in location $B[locator[x]]$
 - ▶ Decrement $locator[x]$

Code for Counting sort

```
def CountingSort(A, B, n , k)
    //Initialize: set each locator[x] to
        the number of entries  $\leq x$ 
    for x = 1 to k do locator[x] = 0
    for i = 1 to n do locator[A[i]] = locator[A[i]] + 1
    for x = 2 to k do
        locator[x] = locator[x] + locator[x-1]
    //Fill output array, updating locator values
    for i = n down to 1 do
        B[locator[A[i]]] = A[i]
        locator[A[i]] = locator[A[i]] - 1
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    for i = n down to 1 do
        B[locator[A[i]]] = A[i]
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```

Analysis: $O(n + k)$ running time.

Counting Sort Example

A:

1	2	3	4	5	6	7	8	9	10
1	3	5	7	5	7	3	8	7	4

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1	2	3	4	5	6	7	8	9	10
1	3	5	7	5	7	3	8	7	4

locator:

1	2	3	4	5	6	7	8
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B:

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			4						

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Counting Sort Example

A:

1	2	3	4	5	6	7	8	9	10
1	3	5	7	5	7	3	8	7	4

locator:

1	2	3	4	5	6	7	8
1	1	1	3	4	6	6	9

B:

1	2	3	4	5	6	7	8	9	10
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B:

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Counting Sort Example

A: Done!

	1	2	3	4	5	6	7	8	9	10
	1	3	5	7	5	7	3	8	7	4

locator:

	1	2	3	4	5	6	7	8
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B:

	1	2	3	4	5	6	7	8	9	10
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- ▶ Simplest approach is to divide the space of possible keys into equal sized buckets.

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 1. **Distribute** keys into buckets
 2. **Sort** keys in each bucket
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- ▶ Simplest approach is to divide the space of possible keys into equal sized buckets.
- ▶ Typically use insertion sort in phase 2.

Bucket Sort Example

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Sort the following keys in the range 0-999, using 10 equal-size buckets:

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1. Distribute

0: 74
1: 140 198 113
2:
3:
4: 467 449
5:
6: 661 642
7:
8: 835
9: 923

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0: 74
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2:
3:
4: 467 449
5:
6: 661 642
7:
8: 835
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2. Sort

0: 74
1: 113 140 198
2:
3:
4: 449 467
5:
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0:	74
1:	140 198 113
2:	
3:	
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5:	
6:	661 642
7:	
8:	835
9:	923

2. Sort

0:	74
1:	113 140 198
2:	
3:	
4:	449 467
5:	
6:	642 661
7:	
8:	835
9:	923

3. Combine

	74
	113
	140
	198
	449
	467
	642
	661
	835
	923

Analysis of Bucket Sort

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n = number of items to sort

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Running time

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Running time

1. Distribution

$O(n)$

Analysis of Bucket Sort

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Phase	Running time
1. Distribution	$O(n)$
2. Sorting each bucket	$O(b + \sum_i s_i^2)$

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- ▶ **Average case:** $O(n)$ if certain assumptions are satisfied (next slide)

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- ▶ **Average case:** $O(n)$ if certain assumptions are satisfied (next slide)
- ▶ **Storage:** is $O(n + b)$.

Average running time of Bucket Sort

The following result is proved in [CLRS]:

Assume:

- 1. The number of buckets is equal to the number of keys (i.e., if $b = n$)*
- 2. The keys are distributed independently and uniformly over the buckets*

Then the expected total cost of the intra-bucket sorts is $O(n)$.

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def radix_sort(A,n);  
  for field ranging from rightmost (least significant)  
    to leftmost (most significant):  
    sort A on field using a stable sort
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Sort the following numbers using radix sort (each digit is a field)

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661

074

835

140

198

923

113

642

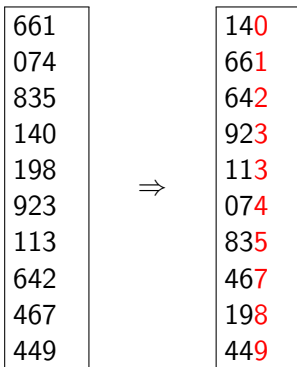
467

449

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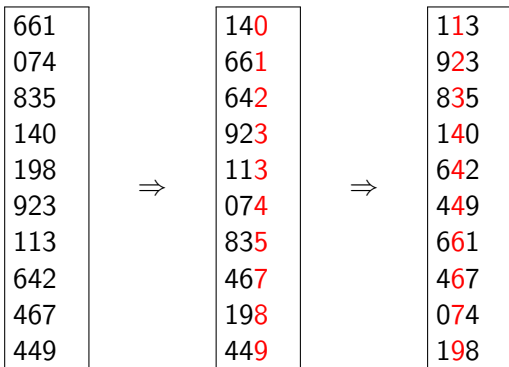
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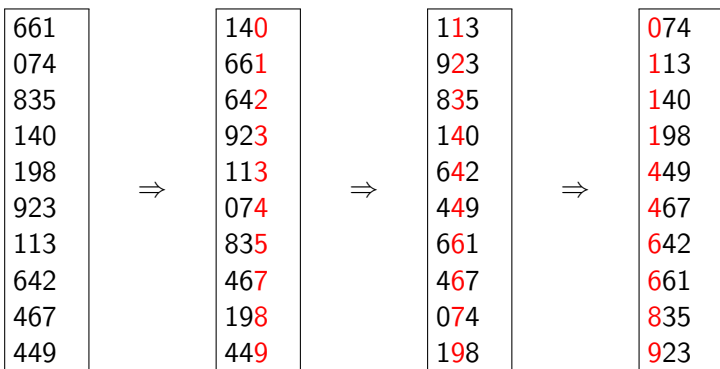
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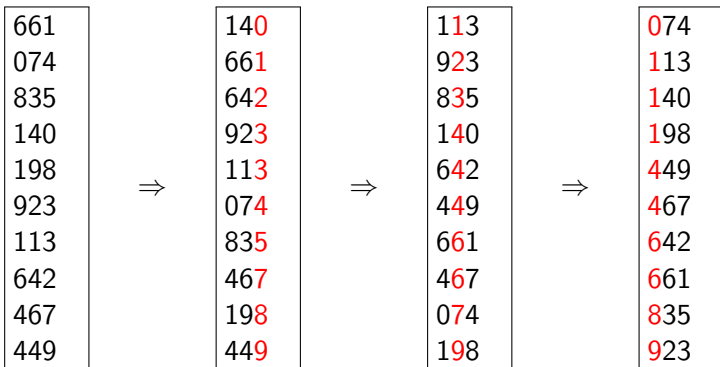
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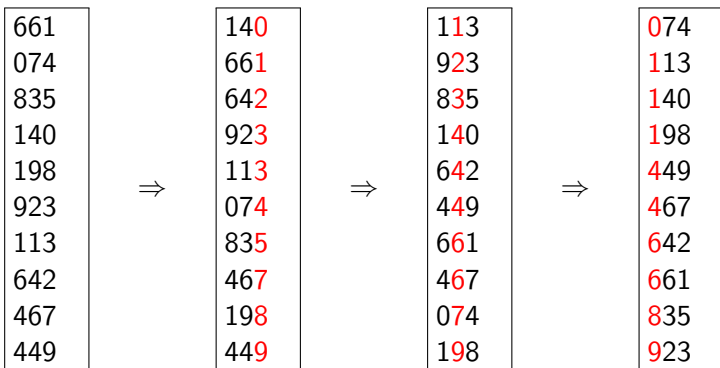


Note the importance of stability.

Radix Sort Example:

Sort the following numbers using radix sort (each digit is a field)

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Note the importance of stability.

We break the ties first, and stability makes sure the ties remain broken correctly.

Analysis of Radix Sort

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- ▶ Each field is sorted using Bucket Sort or Counting Sort

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 - ▶ For example, if each item is a base b number with d digits. (i.e., between 0 and $b^d - 1$, inclusive).
- ▶ Each field is sorted using Bucket Sort or Counting Sort

Then the running time of radix sort is $O(d(n + b))$.

Deterministic Selection: Find k-th element

Recall QuickSelect

quickSelect(S, k)

If n is small, brute force and return.

Pick a random $x \in S$ and put rest into:

L , elements smaller than x

G , elements greater than x

if $k \leq |L|$ **then**

 quickSelect(L, k)

else if $k == |L| + 1$ **then**

return x

else

 quickSelect($G, k - (|L| + 1)$)

Deterministic Selection

Instead of picking x at random:

- ▶ Divide S into $g = \lceil n/5 \rceil$ groups
- ▶ Each group has 5 elements (except maybe g^{th})
- ▶ Find median of each group of 5
- ▶ Find median of those medians
- ▶ Let x be that median.

We call this the “medians of 5” method.

Selecting Median of 5 Example

870	647	845	742	372	882	691	341	461	596
989	151	100	729	101	397	825	587	363	283
595	524	930	259	133	955	620	970	430	280
839	139	735	590	782	913	378	474	255	739
875	150	791	779	792					

Deterministic Select

DeterministicSelect(S, k)

If n is small, brute force and return.

Pick $x \in S$ via medians-of-5 and put rest into:

L , elements smaller than x

G , elements greater than x

if $k \leq |L|$ **then**

DeterministicSelect(L, k)

else if $k == |L| + 1$ **then**

return x

else

DeterministicSelect($G, k - (|L| + 1)$)

Deterministic Selection

Let's visualize: how does pivot compare to list?

Demo Re-visualized

- ▶ Each column was a group of five.
- ▶ Each column is sorted
- ▶ Columns are ordered based on median-of-5
- ▶ Which cells are in L ? G ? Either?

100	283	255	133	341				
101	363	378	259	461				
151	397	474	524	596	620	735	742	791
				691	955	782	845	792
				882	970	839	870	875

Deterministic Selection

- ▶ How few elements *must be* smaller than pivot?
- ▶ How few *must be* non-smaller than pivot?
- ▶ How many could be in either group?